

Graph Subshifts

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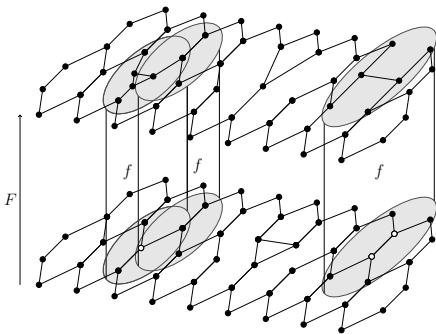
Laboratoire d'Informatique et Systèmes (LIS)

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Causal graph dynamics

Causal graph dynamics (CGD) are an extension of cellular automata:

- The underlying grid becomes an arbitrary graph of bounded degree.
- The graph itself can evolve over time.



However the rule remains *local* and *shift-invariant*.

Group subshifts

Let Σ be a finite alphabet, and G be a finitely generated group.
A **configuration** is a coloring $x : G \rightarrow \Sigma$.

Definition

A **subshift** X is a subset of Σ^G such that there exist a set \mathcal{F} of forbidden **patterns** such that :

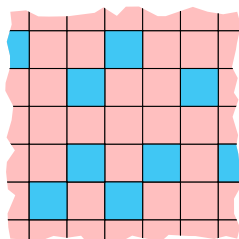
$$X = \{x \in \Sigma^G \mid x \text{ contains no patterns of } \mathcal{F}\}$$

X is said to be of **finite type** if it can be defined by a finite \mathcal{F} .

$$G = \mathbb{Z}^2$$

$$\Sigma = \{ \text{blue square}, \text{red square} \}$$

$$\mathcal{F} = \{ \text{two blue squares horizontally}, \text{two blue squares vertically} \}$$



Subshifts and cellular automata

Proposition (Folklore)

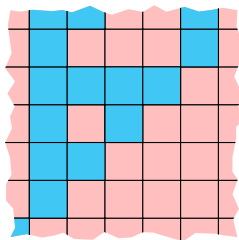
Any cellular automaton of dimension d can be seen as a subshift of finite type in dimension $d + 1$.

For example, the XOR of neighborhood two can be implemented as the subshift :

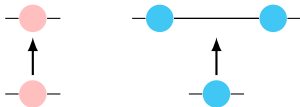
$$G = \mathbb{Z}^2$$

$$\Sigma = \left\{ \begin{array}{|c|} \hline \text{blue} \\ \hline \end{array} \quad \begin{array}{|c|} \hline \text{red} \\ \hline \end{array} \right\}$$

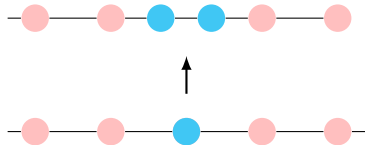
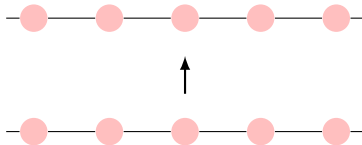
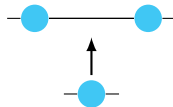
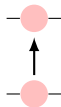
$$\mathcal{F} = \left\{ \begin{array}{|c|c|} \hline & \text{blue} \\ \hline \text{blue} & \text{blue} \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline & \text{red} \\ \hline \text{red} & \text{blue} \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline & \text{red} \\ \hline \text{blue} & \text{red} \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline & \text{blue} \\ \hline \text{red} & \text{red} \\ \hline \end{array} \right\}$$



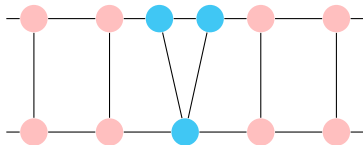
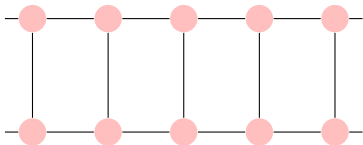
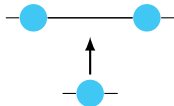
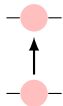
Doesn't work for graph dynamics



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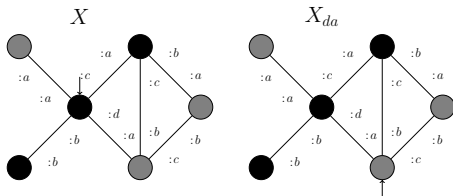
Doesn't work for graph dynamics



Definition (Configuration/Graph)

- A set V of vertices, with a special origin vertex p the pointer
- A set π of ports
- A set E of edges, where an edge is an unordered pair $\{u : a, v : b\}$ with $u, v \in V$ and $a, b \in \pi$. A port is used at most once.
- A coloring function σ from V to Σ .

The graphs are considered modulo isomorphisms. \mathcal{X} is the set of graphs.



Definition (Cuts)

Let $\mathcal{L} \subseteq \pi^*$ be a prefix-stable language. The cut $X|_{\mathcal{L}}$ is the subgraph induced by the vertices reachable through \mathcal{L} .

Definition (Graph Subshift)

Let \mathcal{F} be a set of tuples (F, L) , with F a finite graph and L a prefix-stable language. The subshift forbidding \mathcal{F} is

$$\mathcal{Z} = \{ X \in \mathcal{G}_{\Sigma, \pi} \mid \forall v \in X, \forall (F, L) \in \mathcal{F}, (X_v)|_L \neq F \}.$$

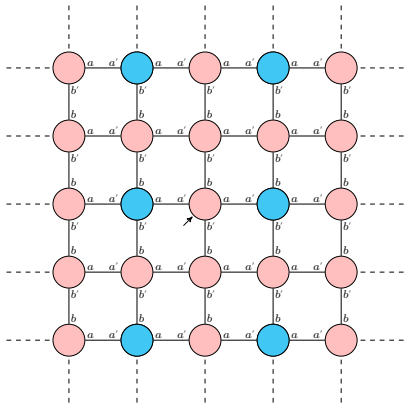
It is said to be of finite type if \mathcal{F} is finite.

Example : Generalized Hardsquare model

$$\Sigma = \{ \text{red circle} \quad \text{blue circle} \}$$

$$\pi = \{a, a', b, b'\}$$

$$\mathcal{F} = \left\{ \begin{array}{l} \text{red} \xrightarrow{a} \text{red} \\ \text{red} \xrightarrow{a'} \text{red} \\ \text{red} \xrightarrow{b} \text{blue} \\ \text{red} \xrightarrow{b'} \text{blue} \\ \text{blue} \xrightarrow{a} \text{red} \\ \text{blue} \xrightarrow{a'} \text{red} \\ \text{blue} \xrightarrow{b} \text{blue} \\ \text{blue} \xrightarrow{b'} \text{blue} \end{array} \right\}$$

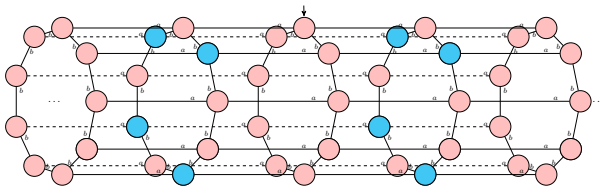


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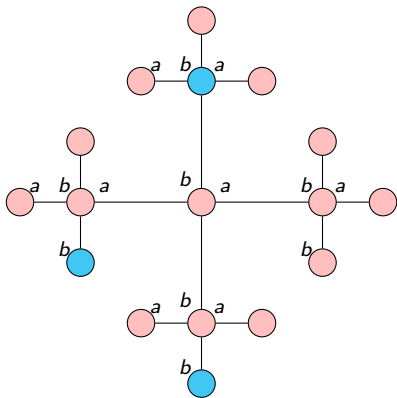


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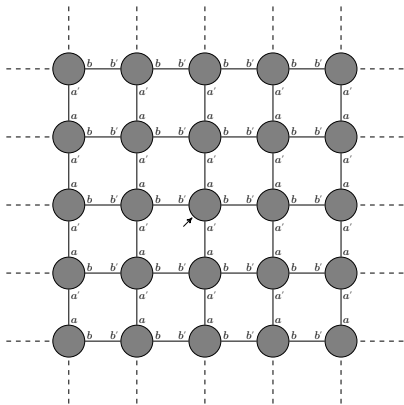
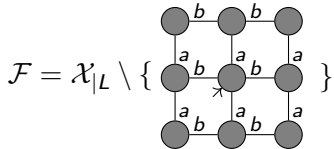
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Example : Locally Grid-like

$$\Sigma = \{ \bullet \}$$

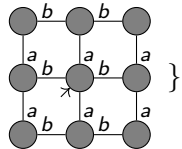
$$\pi = \{a, a', b, b'\}$$

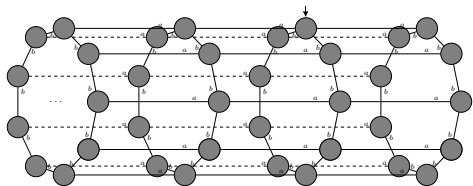


Example : Locally Grid-like

$$\Sigma = \{ \bullet \}$$

$$\pi = \{a, a', b, b'\}$$

$$\mathcal{F} = \mathcal{X}_L \setminus \left\{ \begin{array}{ccc} \bullet & b & \bullet & b & \bullet \\ a & b & a & b & a \\ \bullet & b & \bullet & b & \bullet \\ a & b & a & b & a \end{array} \right\}$$




Definition (Quotient)

Let X be a graph. A X' is a quotient graph of X if there is a surjective homomorphism from X to X' and write $X \succeq X'$.

Definition (Periodicity)

A period of a graph X is a word $w \in \Pi^$ such that $X = X_w$. The stabilizer of X is its set of periods.*

Definition (Strong periodicity)

A graph X is said to be strongly periodic if its stabilizer $\text{STAB}(X)$ is dense in X .

Proposition

The stabilizer of a configuration is a group under the concatenation operation.

Theorem

A graph is strongly periodic if and only if it admits a finite quotient.

Theorem

Let \mathcal{Y} a SFT. \mathcal{Y} admits a strongly periodic configuration X such that $\text{STAB}(X)$ is residually finite if and only if it admits a finite graph.

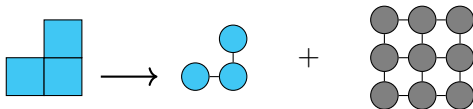
Theorem

The finite configuration problem is undecidable.

Theorem

The support graph unicity problem is undecidable.

Proof idea : Take a \mathbb{Z}^2 subshift and make it a locally \mathbb{Z}^2 -like graph subshift.



To resume: subshifts on things that are not groups.

Perspectives: everything that has already been done in group subshifts.

Thank you for listening !